Graphana - Operations and types

Contents

1 Commands

Every box in this section depicts one operation. The boxes are structured as follows:

operationKey

`parameterName1: ParameterType1`
`parameterName2: ParameterType2`
`...`
`parameterNameN: ParameterTypeN`
`[...]`
returns `ReturnType`

The operation's description text.

Default values are denoted with a `=` after the parameter type followed by the value. If a parameter has a default value then passing an argument is optional. Some operations have three dots at the end of their parameter lists. These operations can receive arguments of the type of the last parameter in the list at any number.

Besides operations, some subsections contain descriptions of terms. These are not written in boxes.

1.1 Program configuration

getCurrentGraph
`deepCopy: Boolean = false` **returns** `Graph`
Returns the current graph. The returned **Graph** can be stored in a variable for example.

setCurrentGraph
`graph: Graph` **returns** `void`
Sets the given graph as the current graph.

setAlgorithmTimeout
`timeOutMillis: PositiveInteger` **returns** `void`
Sets the maximum computation time for an algorithm. If an algorithm which is executed afterwards exceeds the given time then the computation will be aborted and a timeout error will be returned. If 0 is given, then the timeout is disabled. The timeout is given in milliseconds, so a timeout of 1000 means one second. Initially, the timeout is set to 10000.

setScriptTimeout
`timeOutMillis: PositiveInteger` **returns** `void`
Sets the maximum script execution time. If a script which is executed afterwards exceeds the given time then the execution will be aborted and a timeout error will be returned. If 0 is given, then the timeout is disabled. The timeout is given in milliseconds, so a timeout of 1000 means one second. Initially, the timeout is disabled.

setTimeout
`timeOutMillis: PositiveInteger = 10000` **returns** `void`
Sets the algorithm timeout and the script timeout to the given value (see **setAlgorithmTimeout** and **setScriptTimeout**). The value is given in milliseconds.

setPrintWarnings
`printWarnings: Boolean = true` **returns** `void`
Calling this method enables or disables the output of warnings.

setCaching
`enableCaching: Boolean` **returns** `void`
This operation can be used to enable or disable caching of algorithm results. Some algorithms save their (interim) results to reuse them when called repeatedly or to provide them to other algorithms to increase the overall program performance. Initially, caching is enabled. There are circumstances under which caching is automatically disabled, for example if the runtime of an algorithm is measured.

1.2 System operations

import
`Class: ExistingFile` **returns** `String`
Imports the given **ExistingFile** into the program. The file must be a java class which is compatible with *Graphana*. After importing, the operations that are defined within the class are available in the program.

sleep
`milliseconds: PositiveInteger` **returns** `void`
Causes the program to sleep. The duration is given in milliseconds, so for example `1000` means one second.

1.3 Time and date

getTime
`format: String = 'HH:mm:ss'` **returns** `String`
Returns the current system time as formatted **String** in the given format.

millisToString
`milliseconds: PositiveInteger` **returns** `String`
Converts the given `milliseconds` into a formatted **String**.

getTimeMillis
returns `Integer`
Returns the current system time in milliseconds where 0 is 00:00.

getDate
`format: String = 'yyyy/MM/dd'` **returns** `String`
Returns the current system date as formatted **String** in the given format.

1.4 Counters

startCounter
returns `void`
Starts the global counter. Every time this operation is called, the global counter will be reset.

getCounter
returns `Integer`
Returns the time difference between the call of startCounter and the current time in milliseconds. This operation does not stop the counter.

Algorithm timer:
The algorithm timer can be used to measure the runtimes of algorithms. It increases whenever an algorithm is running. So the timer is more accurate than the normal counter because only the runtime of the algorithm itself is measured, ignoring for example compatibility checks. Nevertheless, interferences with the java garbage collector may occur.
The algorithm timer is used via **startAlgorithmTimer** and **getAlgorithmTime**.

startAlgorithmTimer
returns `void`
Starts/restarts the **algorithm timer**. That means, that its value is set to 0.

getAlgorithmTime
returns `Integer`
Returns the current algorithm timer as **Integer** in milliseconds. The algorithm timer keeps running after calling this operation.

1.5 Execution

script

`file: ExistingFile`
`statements: ANY = ''`
`...`
returns `ANY`

Executes the given **ExistingFile** as batch. The script must contain source code in *Graphana* syntax. The additional arguments are ignored and can be used to set up global variables which are used within the script for example.

executeString
`statement: String` **returns** `ANY`
Executes the given **String** and returns the result of the execution. The given string must be source code in *Graphana* syntax. The additional arguments are ignored and can be used to set up global variables which are used within the statement for example. If the statement shall be executed multiple times it is recommended to use **parse** and **executeTree** instead of this command.

parse
`source: String` **returns** `ParseTree`
Parses the given **String** and returns a **ParseTree**. The given `source` must be source code in *Graphana* syntax.

parseScript
`script: ExistingFile` **returns** `ParseTree`
Parses the given **ExistingFile** and returns a **ParseTree**. The script must be source code in *Graphana* syntax.

executeTree
`tree: ParseTree` **returns** `ANY`
Executes the given **ParseTree** and returns the result of the execution. The execution of a parse tree is much faster than the execution of a **String** with **executeString**.

1.6 System alerts

error

`message: String`
returns `void`

Throws an error with the given text message and stops the execution of the statement and, if executed in a script, of the script.

warning

`message: String`
returns `void`

Prints a warning with the given text together with some meta data.

alert

`message: String`
`title: String = 'Message'`
returns `void`

Shows a message dialogue containing the given text. The dialogue window will be titled with the given `title`.

1.7 File output

setOutputFile

`file: File`
`autoWriteConsoleOutput: Boolean = false`
`autoWriteConsoleInput: Boolean = false`
returns `void`

Sets the current output file. After the output file is set, every `WRITE` call will write into the chosen file.
If the given file does not exist, it will be created. Otherwise it will be overwritten. If `autoWriteConsoleOutput` is set to true then nearly every console output, including `PRINT` calls, errors and warnings, will be written into the file automatically. If `autoWriteConsoleInput` is set to true then also console inputs will be written into the file.

flushOutput

returns `void`

Flushes the current output file without closing it. So the file will be visible and up to date in the file system.

closeOutput

returns `void`

Closes the current output file. So the file will be visible and up to date in the file system and can be used by other programs. After closing the file it is not allowed to call `WRITE` until a new output file is set using **setOutputFile**.

writeWholeFile

`file: File`
`object: ANY`
returns `void`

Creates a text file which contains the string representation of the given **Object**. The file will be automatically closed after writing.

1.8 File input

readWholeFile
`file: ExistingFile` **returns** `String`
Reads the whole given (text) file and returns the content as one **String**.

getFiles

`directory: String`
`acceptedExtensions: Vector = )`
returns `Set`

Returns all files of the given directory as a set of **File**. Instead of a directory, a filename can be given alternatively. In this case, a set, which only contains the given file, will be returned.

1.9 Graph visualization

Visualization window:
Every graph visualization and **algorithm visualization** is done in a visualization window which can be minimized, maximized and closed. Within the window, the following actions can be performed:

Left click on a vertex: moving vertex.
Right click on empty space: adding a vertex.
Middle click and drag: scrolling through the view.
Mouse wheel: zoom in and out.
Right click and hold on a vertex and release on another vertex: creating an edge from the first vertex to the second or delete the respective edge, if it already exists.

Modifying the graph only works in the standard visualization and only if it is allowed (for example it is not possible in algorithm visualizations). So the right mouse button may have no effect.

Every window has a certain frame rate which determines, how often the graph is repainted per second. Repainting is necessary to make changes in the dates and states of the vertices and edges visible. If the frame rate is set to zero, then the graph must be repainted manually using **repaintGraph**.

Layout:
The layout determines how the vertices are positioned in a **visualization window**. The layout

is chosen for example as the first argument of the **showGraph** operation.

The following layouts are available in *Graphana*:
GRID
CYCLE
TREE
JUNG.CYCLE
JUNG.ISOM

For directed graphs, a root vertex must be given for every connected component when using the TREE layout. These are specified by writing a colon and the vertex identifiers seperated by commas (for example TREE:v1,v5).

showGraph

`layout: String = 'GRID'`
`windowTitle: String = ''`
`width: PositiveInteger = 640`
`height: PositiveInteger = 640`
`enableModification: Boolean = true`
`frameRate: PositiveInteger = 10`
returns `void`

Visualizes the graph in the **visualization window** with the given `title`. If no such window exists, a new one will be created. The **layout** is set by passing the respective layout keyword (e.g. "Jung.ISOM","TREE" ...). With `width` and `height` the dimensions of the window can be set. If `allowModification` is set to `false` then the graph cannot be modified within the visualization window, so right click will not have any effect. The parameter `frameRate` sets the frame rate of the visualization window. If zero is given then the window will not update frequently.

repaintGraph
`windowTitle: String = 'Graph'` **returns** `void`
Refreshes the graph visualization in the **visualization window** with the given title.

closeGraphView
`windowTitle: String = 'Graph'` **returns** `void`
Closes the graph **visualization window** with the given title.

Algorithm visualization:
Some algorithms support algorithm visualization, which is a step-by-step algorithm output. If algorithm visualization is enabled and a respective algorithm is executed, the visualization starts automatically. Depending on the algorithm, the graph or multiple graphs are visualized in one or more **visualization window**(s) after every important step of the algorithm. One can iterate through the steps by pressing enter in the console. To abort the visualization, 'fin' can be typed in.
The algorithm visualization blocks **caching** and the **algorithm timer**. Algorithm visualiza-

tion can be enabled or disabled using **setAlgorithmOutput**. Initially, algorithm visualization is disabled.

setAlgorithmVisualization
`showOutput: Boolean` **returns** `void`
This operation enables or disables **algorithm visualization**.

setAlgorithmVisualizationParams

`layout: String = 'GRID'`
`width: PositiveInteger = 640`
`height: PositiveInteger = 640`
returns `void`

Sets the visualization parameters for the **algorithm visualization**, which can be enabled using **setAlgorithmOutput**. The parameters have the same meaning as the respective parameters in the **showGraph** operation.

1.10 Histogram creation

newHistogram

`estimatedValues: PositiveInteger = 64`
`xScale: PositiveInteger = 1`
returns `Histogram`

Creates a new empty **Histogram**. The returned **Histogram** can be filled with values using **setHistogramValue** or **incHistogramValue**. Initially, a value is zero. With the parameter `estimatedValues` the initial memory allocation can be set. The capacity is nearly unlimited - `estimatedValues` only has a slight effect on performance.

setHistogramValue

`histogram: Histogram`
`index: PositiveInteger`
`value: Float`
returns `void`

Sets the value associated with the given index in the given histogram.

incHistogramValue

`histogram: Histogram`
`index: Float`
`incValue: Float = 1.0f`
returns `void`

Increments the value associated with the given index in the given **Histogram** by the given `incValue`, which may be negative, too.

csvToHistogram

`csv: String`
`separator: String = ':'`
returns `Histogram`

Converts a CSV string into a **Histogram** which then can be used for example for visualization.

1.11 Histogram visualization

showHistogram

`histogram: Histogram`
`titleKey: String = 'Histogram'`
`clearPrevious: Boolean = true`
`width: Integer = 800`
`height: Integer = 600`
returns `void`

Visualizes the given **Histogram** using the window with the given title. If no such Window is shown, a new one will be created. If `clearPrevious` is set to `false` then previously shown histograms of the window won't be deleted. The dimension of the output window can be set by the parameters `width` and `height`.

addHistogramToView

`histogram: Histogram`
`titleKey: String = 'Histogram'`
returns `void`

Does the same as **showHistogram** with `ClearPrevious` set to false.

setHistogramViewMode

`linesMode: Boolean`
`thickLines: Boolean = true`
`logScale: Boolean = false`
`beginAtZero: Boolean = true`
returns `void`

Configures histogram visualization in general. This will influence every histogram visualization which is done after this operation. If `linesMode` is set to `true` then lines will be drawn instead of bars. With `boldLines` set to true the lines have a width of 3px instead of 1px.

refreshHistogramView
`titleKey: String = 'Histogram'` **returns** `void`
Refreshes the visualization of histograms associated with the given title. This operation must be called after changing **Histogram** values to make the changes visible to the user.

setHistogramViewColors

`color: Color`
`...`
returns `void`

Configures the colors of the bars or lines of all histogram visualization which is called after this operation. The first given **Color** is used for the first added **Histogram** of a visualization, the second **Color** for the second one and so on. If there are more histograms to output than colors given then it restarts with the first **Color**.

clearHistogramView

`titleKey: String = 'Histogram'`
returns `void`

Removes all the histograms of a visualization associated with the given title. The visualization window remains visible.

getHistogramFromView

`titleKey: String = 'Histogram'`
`index: PositiveInteger = 0`
returns `Histogram`

Extracts the **Histogram** with the given index from a visualization associated with the given title. The indices of the histograms are set by the order they were added into the visualization.

1.12 Colors

newColor

`red: PositiveInteger`
`green: PositiveInteger`
`blue: PositiveInteger`
`alpha: PositiveInteger = 255`
returns `Color`

newFloatColor

`red: Float`
`green: Float`
`Blue: Float`
`Alpha: Float = 1.0f`
returns `Color`

gray

`value: PositiveInteger`
returns `Color`

Returns a gray color with the given brightness. The brightness must be a value between 0 (black) and 255 (white)

fGray
`value: Float` **returns** `Color`
Returns a gray color with the given brightness. The brightness must be a value between 0 (black) and 1 (white)

1.13 User interactions

ask
`question: String = ''` **returns** `String`
Pauses the execution, waits for a user input and returns the **String** which was entered by the user.

pause
`message: String = 'Press Enter...'` **returns** `void`
Pauses the execution until the user presses enter. A message can be given. This message will be printed before the execution pauses.

1.14 Variables

typeOf
`variable: ANY` **returns** `String`
Returns the type name of the given value.

removeVariable
`identifier: String` **returns** `Boolean`
Removes the variable with the given `identifier`. The value will be deleted from memory and calling `defined` on the variable afterwards will return `false`.

1.15 Assertions

assert

`condition: Boolean`
`message: ANY = ''`
returns `void`

Does nothing, if the given **Boolean** is true. Otherwise, an error is thrown together with a message that can be given.

assertEq

`value1: ANY`
`value2: ANY`
`message: ANY = ''`
returns `void`

Does nothing, if the two given values are equal. Otherwise, an error is thrown together with a message that can be given.

1.16 Bounds

newInterval

`lowerBound: Float`
`upperBound: Float`
returns `Interval`

Creates and returns a new **Interval** with the given bounds.

getLowerBound
`bounds: Interval` **returns** `Float`
Returns the lower bound of the given **Interval**.

getUpperBound
`bounds: Interval` **returns** `Float`
Returns the upper bound of the given **Interval**.

1.17 Converting primitives

asFloat
`integer: Integer` **returns** `Float`
Converts an **Integer** into a **Float**. This can be used for example to enforce float division when dividing two integers.

parseInt
`string: String` **returns** `Integer`
Converts a **String** into an **Integer** by parsing the string.

parseFloat
`string: String` **returns** `Float`
Converts a **String** into a **Float** by parsing the string.

parseBool
`string: String` **returns** `Boolean`
Converts a **String** into a **Boolean** by parsing the string. The strings `''true''` and `''1''` result in `true` and the strings `''false''` and `''0''` result in `false`. The strings are not case-sensitive.

1.18 String operations

toString
`object: ANY` **returns** `String`
Returns the **String** representation of the given `object` which can be of any type.

split

`string: String`
`regex: String = '\n'`
`trim: Boolean = true`
returns `Vector`

Splits the given **String** at the given regular expression and returns multiple strings as a **Vector**.

startsWith

`string: String`
`prefix: String`
returns `Boolean`

Returns `true` iff the given **String** starts with `prefix`.

endsWith

`string: String`
`postfix: String`
returns `Boolean`

Returns `true` iff the given **String** ends with `postfix`.

1.19 Complex type operations

getSize
`iterable: Iterable` **returns** `PositiveInteger`
Returns the number of elements in the given **Iterable**.

getVectorSize
`vector: Vector` **returns** `PositiveInteger`
Returns the number of entries of the given **Vector**.

setVectorSize

`vector: Vector`
`newSize: PositiveInteger`
returns `void`

Sets the number of entries of the given **Vector** to the given number. The values of the vector remain the same.

getSetCardinality
`set: Set` **returns** `PositiveInteger`
Returns the cardinality of the given **Set**.

setInsert

`set: Set`
`value: ANY`
returns `void`

Inserts the given element into the given **Set**. The element is inserted even if an equal element exists in the given set.

1.20 Math functions

round
`number: Float` **returns** `Integer`
Converts a **Float** into an **Integer** by rounding the given value.

random

`lowerBound: Integer`
`upperBound: Integer`
returns `Integer`

Returns a random **Integer** which is bigger or equal to `lowerBound` and smaller or equal to `upperBound`.

sqrt
`x: Float` **returns** `Float`
Returns the square root of the given value.

sqr
`x: Float` **returns** `Float`
Returns the square of the given value.

pow

`base: Float`
`exp: Float`
returns `Float`

Returns `base` to the power of `exp`.

sin
`x: Float` **returns** `Float`
Returns the sine of the given value.

cos
`x: Float` **returns** `Float`
Returns the cosine of the given value.

tan
`x: Float` **returns** `Float`
Returns the tangent of the given value.

cotan
`x: Float` **returns** `Float`
Returns the cotangent of the given value.

2 Graph operations

In the explanations of this section, G is the given graph, V its vertices and E its edges. When a runtime is given then n is the number of vertices, m the number of edges and $\Delta(G)$ is the sum of both.

Every box in this section depicts one algorithm. The boxes are structured as follows:

algorithmKey

```
parameterName1:  ParameterType1
parameterName2:  ParameterType2
...
parameterNameN: ParameterTypeN
[...]
returns ReturnType
```

The algorithm's description text.

For some algorithms:
Runtime: The algorithms runtime in O-notation
Graph preconditions: List of preconditions
Compatible libraries: List of supported graph libraries

Arguments are handled in the same way as they were explained in the previous section. The **algorithm timer** only counts if one of the algorithms of this section is called.

Algorithms which support **algorithm visualization** are marked with a * after the algorithm key.

2.1 Graph creation

Graph configuration:
In *Graphana* a graph configuration is the combination of the properties *directed*, *weighted* and *loop-free forced*. If a graph is forced to be loop-free then no loops can be inserted.

createGraph

```
directed:  Boolean = false
weighted:  Boolean = false
forceLoopFree:  Boolean = false
library:  String = 'KEEP'
returns void
```

Creates a graph with the given **graph configuration** and sets it as the current graph. With the parameter `library` a name of a **graph library** can be given. The graph will then internally be created as a graph of the respective library. If the argument is set to `KEEP` or omitted then the previously used library will be used. An already created graph will be completely deleted and recreated.

2.2 General graph properties

vertexCount
returns `Integer`
Returns the number of vertices.

edgeCount
returns `Integer`
Returns the number of edges.

graphSize
returns `Integer`
Returns the sum of the vertex count and edge count.

isDirected
returns `Boolean`
Returns `true`, if and only if the graph is directed.

isWeighted
returns `Boolean`
Returns `true`, if and only if the graph is weighted.

isLoopFree
returns `Boolean`
Returns `true`, if and only if the graph does not contain any loops.

isForceLoopFree
returns `Boolean`
Returns `true`, if no loops can be added to the graph.

2.3 Graph loading

loadGraph
`filename: ExistingFile` **returns** `void`
Sets the current graph by loading a DIMACS or a dot file. Depending on the given file format, the operation does either the same as **loadDIMACS** or **loadDot**.

loadDIMACS
`filename: ExistingFile` **returns** `void`
Loads the given DIMACS **File**. If the graph is directed then every edge in the file is seen as an directed edge and vice versa. So if the graph is undirected, there can only be one edge per vertex pair even if there are two in the file. If the graph is unweighted then the weights within the file will be ignored. If the graph is forced to be loop-free then loops in the file will be ignored . For huge files the number of lines to read can be limited using 'MaxLines'.

loadDot

`filename: ExistingFile`
`ignoreWeights: Boolean = false`
`ignoreLayoutAttributes: Boolean = false`
returns `void`

Loads the given dot **File**. Since a dot file directly contains information of whether the graph is directed or not, the resulting graph will be directed if and only if it is directed in the dot file. If the graph is unweighted then the weights within the file will be ignored. If the graph is forced to be loop-free then loops in the file will be ignored .

2.4 Graph libraries

Graph library:
Graphana can internally use different graph libraries. Which library is used influences performance and the set of available algorithms. The usage itself does not depend on the chosen library. So graph construction, graph loading etc. always works in the same way. In addition libraries can be converted into each other (either manually by calling **setLibrary** or automatically if a called algorithm is not compatible with the current library).

setLibrary
`libraryName: String` **returns** `void`
Sets the current **graph library**. The graph will be converted into the given graph library. Initially, the JUNG2 library is set.

getLibrary
returns `String`
Returns the name of the current graph library as a **String**.

getAvailableLibraries
returns `Set`
Returns the names of all available graph libraries as a **Set** of **String**. Each of the given names is a valid library input for the **setLibrary** operation or the **createGraph** operation.

2.5 Random graphs

createErdosRenyiGraph

```
vertexAmount: PositiveInteger = 10
connectionProbability: Float = 0.3f
minWeight: Integer = 1
maxWeight: Integer = 1
```
returns `void`

Creates a random Erdős-Rï¿½nyi-Graph. The new graph has `VertexCount` vertices. Every vertex pair is connected with a probability of `ConnectionProbability`. The weight of every edge is a random value between `MinWeight` and `MaxWeight`. These values only have an effect if the graph was previously created as a weighted graph (see **createGraph**).

createRandomClusterGraph

```
clusterAmount: PositiveInteger = 10
minClusterSize: PositiveInteger = 10
maxClusterSize: PositiveInteger = 20
additionalEdgesAmount: PositiveInteger = 8
minWeight: Integer = 1
maxWeight: Integer = 1
```
returns `void`

Creates a random cluster graph. The resulting graph contains up to `ClusterAmount` clusters.

addRandomClique

```
membershipProbability: Float
```
returns `void`

Adds a clique in the graph by adding edges using the existing vertices. Every vertex of the graph is part of the clique with a probability of `MemberShipProbability`. Setting the value to 1 means that the whole graph will be a clique or will be complete, respectively. By setting it to 0, the operation has no effect.

createPGeneratedGraph

```
vertexCount: Integer
a: Float = 0.5f
b: Float = 0.6f
```
returns `void`

Creates a p-generated random graph.
Note that the resulting graph is always undirected, unweighted and loop-free.

2.6 Graph editing

resolveVertexNameClashes

```
resolve: Boolean
```
returns `void`

If `resolve` is `true` then name clashes will be automatically resolved when adding a vertex (e.g. with **addVertex**) with an identifier which is already used by an existing vertex of the graph. Initially name clashes are not resolved.

addVertex

`identifier: String = ''`
returns `Vertex`

Adds a vertex with the given name to the current graph. The added vertex is then identified by the given `identifier`. If auto-resolving of name clashes is activated (see **resolveVertexNameClashes**) then underscores will be added to the given identifier until there is no vertex with the same identifier. If not and there is a name clash then no vertex will be added. If no identifier is given, a default identifier will be used (default identifiers are enumerated). The new or the already existing vertex is returned.

addVertices

`identifier: String`
`...`
returns `void`

Adds multiple vertices. With every given `identifier` the operation adds a vertex just as **addVertex** does

addVertexRow

`amount: PositiveInteger`
`startIndex: Integer = 0`
`prefix: String = 'v'`
`Cluster: Boolean = false`
returns `void`

Adds overall `amount` vertices. The operation enumerates the added vertices, starting at `startIndex`. The name of an added vertex will be the `prefix` concatenated with the number. If `cluster` is set to `true` then all added vertices are connected with each other.

addEdge

`startVertex: Vertex`
`endVertex: Vertex`
`weight: Float = 1.0f`
returns `void`

Adds an edge between the two given vertices (see **Vertex**). A weight can be given, but will be ignored, if the graph is unweighted.

setEdgeWeight

`edge: Edge`
`weight: Float = 1.0f`
returns `void`

Sets the weight of the given **Edge**. An error is returned if the graph is unweighted.

removeVertex

`vertex: Vertex`
`...`
returns `void`

Removes the given **Vertex** or the given vertices, respectively, from the graph. That means one ore more vertices can be given.

removeVertexSet

`vertices: Iterable`
`...`
returns `void`

Removes all vertices of the given **Iterable**.

removeEdge

`edge: Edge`
`...`
returns `void`

Removes the given **Edge** or the given edges, respectively, from the graph. That means one or more edges can be given.

removeEdgeSet

`edges: Iterable`
`...`
returns `void`

Removes all edges of the given **Iterable**.

clearGraph
returns `void`
Removes all vertices from the graph.

deleteLoops
returns `void`
Deletes all loops from the graph in order that the graph is loop-free after this operation. However, loops can be inserted afterwards. To disallow this, see **forceLoopFree**.

forceLoopFree
returns `void`
Deletes all loops from the graph. Furthermore, loops cannot be inserted afterwards.

allowLoops
returns `void`
After the call of this operation, loops can be added into the graph.

mergeGraph
`sourceGraph: Graph` **returns** `void`
Merges the graph with the given `sourceGraph`. Every vertex and edge of the given graph will be added to the graph as deep copies. Only dates of the vertices and edges, if existing, are not copied deep.

graphGUI

`deleteGraph: Boolean = false`
`drawWindowWidth: PositiveInteger = 640`
`drawWindowHeight: PositiveInteger = 640`
`frameRate: PositiveInteger = 0`
returns `void`

Opens a **visualization window** with the standard grid layout with the purpose of editing the graph visually. If `deleteGraph` is set to `true` then all vertices are deleted before editing and the graph as well as the visualization window is empty.

2.7 Graph conversions

setGraphConfig

`directed: Boolean`
`weighted: Boolean`
`forceLoopFree: Boolean`
returns `void`

Converts the current graph into a graph with the given **graph configuration** whereas the graph library remains the same. If the given graph configuration is forbidden in the respective graph library then an error will be returned.

asDirected

returns `Graph`

Returns an equivalent directed graph. The returned graph contains the same vertices as the original graph. For every undirected edge in the original graph two directed edges are created in the returned graph. If the original graph is already directed then the graph is returned without any changes.

toDirected

returns `void`

Converts the current graph into a directed graph. The converted graph contains the same vertices as the original graph. For every undirected edge in the original graph two directed edges are created in the converted graph.
If the current graph is already directed then nothing happens.

asWeighted

returns `Graph`

Returns an equivalent weighted graph. The returned graph contains the same vertices as the original graph. For every unweighted edge of the original graph, an edge with the weight 1 is created in the returned graph. In the returned graph, edge weights can be set.
If the original graph is already weighted then the graph is returned without any changes.

toWeighted
returns `void`
Converts the current graph into a weighted graph. The converted graph contains the same vertices as the original graph. For every unweighted edge of the original graph an edge with the weight 1 is created in the converted graph. After this call, edge weights can be set in the current graph. If the current graph is already weighted then nothing happens.

graphAsDIMACS
returns `String`
Returns a **String** containing the DIMACS representation of the graph.

3 Algorithms

Algorithms are special Graph operations. The boxes in this section are structured the same way as in the previous section.

Algorithms which support **algorithm visualization** are marked with a * after the algorithm key.

3.1 Vertex degrees

averageDegree
returns `Float`
Returns the average degree of all vertices. **Graph preconditions:** not empty

maxDegree
returns `Integer`
If the graph is undirected then the degree of the vertices with the largest number of neighbors is returned. Otherwise the maximum of **maxIngoingDegree** and **maxOutgoingDegree** is returned. **Graph preconditions:** not empty

maxIngoingDegree
returns `Integer`
Returns the ingoing edge count of the vertices with the largest number of ingoing edges. If the graph is undirected then the returned value is equal to the **maxDegree** return value. **Graph preconditions:** not empty

maxOutgoingDegree
returns `Integer`
Returns the outgoing edge count of the vertices with the largest number of outgoing edges. If the graph is undirected then the returned value is equal to the **maxDegree** return value. **Graph preconditions:** not empty

minDegree
returns `Integer`
If the graph is undirected then the degree of the vertices with the smallest number of neighbors is returned. Otherwise the minimum of **minIngoingDegree** and **minOutgoingDegree** is returned. **Graph preconditions:** not empty

minIngoingDegree

returns `Integer`

Returns the ingoing edge count of the vertices with the smallest number of ingoing edges.
If the graph is undirected then the returned value is equal to the **minDegree** return value.

Graph preconditions: not empty

minOutgoingDegree

returns `Integer`

Returns the outgoing edge count of the vertices with the smallest number of outgoing edges.
If the graph is undirected then the returned value is equal to the **minDegree** return value.

Graph preconditions: not empty

degreeDistribution

returns `Histogram`

Returns a **Histogram** with a mapping from vertex degrees to the amount of vertices that have the respective degree.

Graph preconditions: undirected, not empty

ingoingDegreeDistribution

returns `Histogram`

Returns a **Histogram** with a mapping from vertex degrees to the amount of vertices that have the respective ingoing degree.
If the graph is undirected then the returned histogram is equal to the **degreeDistribution** return value.

Graph preconditions: not empty

outgoingDegreeDistribution

returns `Histogram`

Returns a **Histogram** with a mapping from vertex degrees to the amount of vertices that have the respective outgoing degree.
If the graph is undirected then the returned histogram is equal to the **degreeDistribution** return value.

Graph preconditions: not empty

distanceDistribution

returns `Histogram`

Returns a mapping of d to the number of vertices that have the distance d.

Graph preconditions: not empty
Compatible libraries: JUNG2

3.2 Flows

maxFlow

`source: Vertex`
`sink: Vertex`
returns `Float`

Returns the max flow between the two given vertices (see **Vertex**). **Graph preconditions:** not empty **Compatible libraries:** JUNG2, JGraphT

minCut

`source: Vertex`
`sink: Vertex`
returns `Vector`

Returns the min cut between the two given vertices (see **Vertex**). **Graph preconditions:** not empty **Compatible libraries:** JUNG2

Gomory-Hu-Tree:
The *Gomory-Hu-Tree* $T = (V, E_T)$ of a graph $G = (V, E_G)$ is a tree in order that every pair $(v, w) \in V$ has the same max flow as in G.

gomoryHuTree *
`ignoreWeights: Boolean = false` **returns** `Graph`
Returns the **Gomory-Hu-tree** of the graph. **Runtime:** $O(n^2 + m^2)$ **Graph preconditions:** undirected, not empty, loop-free **Compatible libraries:** JUNG2

3.3 Connected Components

getConnectedComponentCount *
returns `PositiveInteger`
Returns the number of connected components.

getConnectedComponent *
`componentIndex: PositiveInteger` **returns** `Graph`
Returns the connected component with the given index. The index must be a value between 0 and the connected component count (see getConnectedComponentCount) minus one. The indices are given internally. The returned graph is a deep copy of the respective connected component. **Graph preconditions:** undirected, not empty

getConnectedComponentByVertex *
`vertex: String` **returns** `Graph`
Returns the connected component in which the given vertex is contained. The returned graph is a deep copy of the respective connected component. **Graph preconditions:** undirected, not empty

getStronglyConnectedComponentCount
returns `Integer`
Returns the number of strongly connected components. **Graph preconditions:** loop-free

3.4 Trees

Tree:
A *tree* is an acyclic graph.

isTree *
returns `Boolean`
Checks, whether the graph is a **tree**. **Graph preconditions:** undirected

3.5 Treewidth

setTreewidthUpperBoundHeuristics

`heuristic: String`
`...`
returns `void`

Since the computation of the treewidth is NP-complete, *Graphana* uses some heuristics for this problem. The heuristics are implemented in `LibTW` from `www.treewidth.com` .
The heuristics which are to be used are passed as **String**s. If multiple heuristics are given then every heuristic will be executed and the best result will be returned .

The following **String**s are valid treewidth upper bound heuristic keys:
`GREEDYFILLIN`
`GREEDYDEGREE`
`ALLSTARTLEXBFS`

By default, `GREEDYFILLIN` is set. For informations on the different heuristics, see `www.treewidth.com` .

The chosen treewidth upper bound heuristics influence the following algorithms: **treewidthUpperBound**, **treewidthBounds**, **treewidthExact**.

setTreewidthLowerBoundHeuristics

`heuristic: String`
`...`
returns `void`

Since the computation of the treewidth is NP-complete, *Graphana* uses some heuristics for this problem. The heuristics are implemented in `LibTW` from `www.treewidth.com` .
The heuristics which are to be used are passed as **String**s. If multiple heuristics are given then every heuristic will be executed and the best result will be returned .

The following **String**s are valid treewidth lower bound heuristic keys:
`MAXMINDEGREEPLUSLEASTC`
`MAXCARDSEARCH`
`RAMACHANDRAMURTHI`
`ALLSTARTMAXCARDSEARCH`
`MAXMINDEGREE`
`MAXMINDEGREEPLUSMAXD`
`MAXMINDEGREEPLUSMIND`
`ALLSTARTMAXMINDEGREE`
`ALLSTARTMAXMINDEGREEPLUSLEASTC`
`ALLSTARTMINORMINWIDTH`
`MINORMINWIDTH`
`MINDEGREE`

By default, `MAXMINDEGREEPLUSLEASTC` is set. For informations on the different heuristics, see `www.treewidth.com` .

The chosen treewidth lower bound heuristics influence the following algorithms: **treewidthLowerBound**, **treewidthBounds**.

treewidthUpperBound

returns `Integer`

Returns an upper bound of the treewidth. The heuristics which are to be used can be set with **setTreewidthUpperBoundHeuristics**.

Graph preconditions: not empty
Compatible libraries: LibTW

treewidthLowerBound

returns `Integer`

Returns a lower bound of the treewidth. The heuristics which are to be used can be set with **setTreewidthLowerBoundHeuristics**.

Graph preconditions: not empty
Compatible libraries: LibTW

treewidthExact

returns `Integer`

Returns the treewidth using the 'TreewidthDP' algorithm from `www.treewidth.com`. The algorithm has a NP runtime. Before the actual computation starts, an upper bound is established by using one or more heuristics. Which heuristics are to be used for this can be set with **setTreewidthUpperBoundHeuristics**.
For further information on the algorithm, see `www.treewidth.com`.

Graph preconditions: not empty
Compatible libraries: LibTW

3.6 Connectivity

largestKConnected *

`k: Integer`
returns `Integer`

Returns the cardinality of a maximum $V' \subset V$ in order that V' is k-edge-connected depending on the parameter `k`.

Graph preconditions: undirected, not empty, loop-free
Compatible libraries: JUNG2

edgeConnectivityDistribution *

returns `Histogram`

Returns a mapping from k to **largestKConnected**(k) as a **Histogram**. The first value is $k = 0$. The last value is the largest k where **largestKConnected**(k) does not return `0`.

Graph preconditions: not empty
Compatible libraries: JUNG2

3.7 Clusters

Cluster:
A *cluster* is a connected component in which all vertices are connected with each other.

Cluster graph:
A *cluster graph* is a graph which consists only of clusters (see **Cluster**).

isClusterGraph *
returns `Boolean`
Returns `true` if and only if the graph is a **Cluster graph**. **Graph preconditions:** undirected

3.8 Cluster vertex deletion

CVD:
Abbrevation for "cluster vertex deletion"

Cluster vertex deletion set:
A set $C \subseteq V$ in order that $(V \backslash C, E_C)$ is a **cluster graph** (the set of edges $E_C \subseteq E$ contains all edges which are not incident to any vertex in C).

CVD-set:
Abbrevation for **Cluster vertex deletion set**

CVD-heuristics:
Since finding a **CVD-set** is NP-complete, *Graphana* supports several heuristics for this problem which differ in runtime and cardinality of the resulting set.
Which heuristic(s) shall be used, can be set with **setCVDHeuristics**. A CVD-heuristic consists of two parts: the search strategy to search for nodes which may be deleted and the delete strategy to delete one or more vertex of the found candidates. In *Graphana* there are two search strategies and three delete strategies. So in combination, there are six heuristics.

The search strategies are:
Successive (key: "SUCC"):
The candidates are found by regarding each vertexes neighbors. This strategy is recommended for very sparse graphs.
Runtime: $O(n \cdot m)$**Connected components (key: "CC"):**
The candidates are found by recursively splitting the graph into connected components. This strategy is especially recommended for dense graphs. In most cases the runtime is better than the runtime of the "Successive" strategy.

Runtime: $O((n+m) + |C| \cdot (n+m) \cdot \Delta(G) + |C| \cdot (\Delta(G))^2)$
Where $|C|$ is the cardinality of the CVD set.

The delete strategies are:
All (key: "ALL"): Deletes all found candidates. This strategy ensures, that the cardinality of the resulting CVD-set is not more than three times as large as the cardinality of an optimal solution.
First (key: "FIRST"):
Deletes the candidate which was found first.
Maximum degree (key: "MAX")
Deletes a candidate with the highest **degree**.

setCvdHeuristics

`heuristic: String`
`...`
returns `void`

Sets the **CVD-heuristics** which shall be used when computing a **CVD-set**. The heuristics are given as **String**s containing the key of search strategy and the key of the deletion strategy, separated by a minus character. So for example "CC-MAX" is a valid string. More than one heuristic can be set by passing them within one call. A new call of setCVDHeuristics resets the heuristics. If multiple heuristics are given, then the respective algorithms will execute all of them and return the best result. So the computation time increases but the results are getting more accurate. The chosen cvd heuristics influence the following algorithms: **cvdSet**, **cvdSize**, **cvdBounds**, **toClusterGraph**, **maximumIndependentSetByCVD**

toClusterGraph
returns `void`
Deletes vertices in order that the graph becomes a **cluster graph**. The number of deleted vertices may not be optimal (see **CVD-heuristics**). **Graph preconditions:** undirected, not empty, loop-free

getMaximumIndependentSetByCVD
returns `List`
Computes the maximum independent set with a parameterized algorithm which uses a **CVD-set** as parameter. **Graph preconditions:** undirected, not empty, loop-free

getCvdSet
returns `List`
Tries different heuristics to compute a **CVD-set**. **Graph preconditions:** undirected, not empty, loop-free

getCvdSetSize
returns `Integer`
Same as **getCVDSet** but only returns the size of the solution. **Graph preconditions:** undirected, not empty, loop-free

3.9 Cluster Editing

CE:
Abbreviation for "cluster editing". Task: Find a minimum size set of edges to add or delete such that the result is a **cluster graph**.

edge deletion:
The minimum size set of edges to delete such that the result is a **cluster graph**.

clusterEditing
returns `List`
Returns a minimum size set of edges to add or delete such that the result is a **cluster graph**. The edge branching algorithm in use has a runtime of **Runtime:** $O(n^{2.61})$ **Graph preconditions:** undirected, not empty

clusterEditingSize
returns `Integer`
Returns the solution size, see **clusterEditing**. **Graph preconditions:** undirected, not empty

clusterEditing3APX
returns `List`
It is an expected (randomized) 3-approximation for cluster editing. The result is the best of 5 runs. **Graph preconditions:** undirected, not empty

clusterEditing3APXSize
returns `Integer`
Returns the solution size, see **clusterEditing3APX**. **Graph preconditions:** undirected, not empty

clusterEditing3APXRuns
`runs: Integer` **returns** `List`
Same as **clusterEditing3APX**; takes the number of runs as its argument. **Graph preconditions:** undirected, not empty

clusterEditing3APXRunsSize
`runs: Integer` **returns** `Integer`
Returns the solution size, see **clusterEditing3APXRuns**. **Graph preconditions:** undirected, not empty

edgeDeletionAPX
returns `List`
A variaton of **clusterEditing3APX** for **edge deletion**. The result is the best of 5 runs. **Graph preconditions:** undirected, not empty

edgeDeletionAPXSize
returns `Integer`
Returns the solution size, see **edgeDeletionAPX**. **Graph preconditions:** undirected, not empty

edgeDeletionAPXRuns
`runs: Integer` **returns** `List`
Same as **edgeDeletionAPX**; takes the number of runs as its argument. **Graph preconditions:** undirected, not empty

edgeDeletionAPXRunsSize
`runs: Integer` **returns** `Integer`
Returns the solution size, see **edgeDeletionAPXRuns**. **Graph preconditions:** undirected, not empty

3.10 Feedback Sets

Feedback edge set size:
The *feedback edge set size* is the minimum number of edge deletions that are necessary in order that the graph becomes a **tree**.
The feedback edge set size for a connected component is $|E| + 1 - |V|$.

feedbackEdgeSetSize *
returns `Integer`
Returns the **feedback edge set size**. **Graph preconditions:** undirected

feedbackVertexSet *
returns `Set`
 Graph preconditions: undirected, loop-free

feedbackVertexSetSize *
returns `PositiveInteger`
 Graph preconditions: undirected, loop-free

feedbackArcSet
returns `List`
Heuristic for computing the *feedback arc set* of a directed graph; that is the number of edges that have to be removed in order to obtain an acyclic graph. Uses sifting, which is similar to a two sided insertion sort. **Runtime:** $O(n^2)$. **Graph preconditions:** loop-free

feedbackArcSetSize
returns `Integer`
A call to this algorithm is equivalent to **getSize**(**feedbackArcSet**()). **Graph preconditions:** loop-free

3.11 Miscellaneous graph parameters

vertex cover:
A *vertex cover* S is a set of vertices in order that every edge $e \in E$ has at least one endpoint in V. The vertex cover size is the cardinality of a minimum vertex cover.

dominating set:
A *dominating set* S is a set of vertices such that every node in the graph has at least one neighbour in S.

betweenness:
The *betweenness* of a Vertex/Edge is the number of shortest paths that use the Vertex/Edge.

vertexCoverSize
`useGreedy: Boolean = true` **returns** `Integer`
This algorithm offers two heuristics: If `useGreedy` is set to `true` then a *nlogn* - approximation is used. Otherwise a 2-approximation is used. The *nlogn* - approximation delivers better results in many practical cases. **Runtime:** $O(n + m)$ **Graph preconditions:** undirected, loop-free

vertexCoverSizeBothHeuristics
returns `Integer`
Calls both heuristics of **vertexCoverSize** and returns the minimum of both results. **Runtime:** $O(n + m)$ **Graph preconditions:** undirected, loop-free

hIndex
returns `Integer`
The *h-index* is the largest number n in order that n nodes have at least n neighbors. **Runtime:** $O(n + \Delta(G))$ **Graph preconditions:** not empty

hIndexPlus
returns `Integer`
The *h-indexPlus* is the number n of nodes to delete such that the remaining nodes have a degree of at most n. **Graph preconditions:** undirected, loop-free

k-degenerate:
A graph is *k-degenerate* if and only if there is an induced subgraph which contains a vertex with a degree at most k.

degeneracy
returns `Integer`
The *degeneracy* is the smallest number k in order that the graph is **k-degenerate**. **Runtime:** $O(m)$ **Graph preconditions:** undirected, not empty, loop-free

dominatingSet

`useReduction: Boolean = false`
returns `Set`

A heuristic to calculate a (hopefully minimal) dominating set. The result set is returned. The 1-degree reduction rule is always applied. If `useReduction` is set to `true` there will be a further reduction step beforehand (
Runtime: $O(n^3)$).
Graph preconditions: undirected, loop-free

dominatingSetSize

`useReduction: Boolean = false`
returns `Integer`

Same as **dominatingSet** but returns only the size of the result set.

Graph preconditions: undirected, loop-free

betweennessVertex

returns `Histogram`

Calculates the **betweenness** of all vertices and returns a histogram.

Graph preconditions: not empty
Compatible libraries: JUNG2

betweennessEdge

returns `Histogram`

Calculates the **betweenness** of all edges and returns a histogram.

Graph preconditions: not empty
Compatible libraries: JUNG2

neighborhoodDiversity

returns `Integer`

A graph has a neighborhood diversity of at most w, if there exist a partition of V into at most w sets, such that all the vertices in each set have the same neighborhood.

Graph preconditions: not empty

4 Types

Every box in this section depicts one type. The boxes are structured as follows:

TypeName

The type's description text.

Samples:
sample1
sample2
...
sampleN

4.1 Primitive types

Integer

An integral number with the range -2,147,483,648 to 2,147,483,647.

Samples:
8
-10
0

PositiveInteger

Essentially the same as **Integer** but with the range 0 to 2,147,483,647.

Samples:
3
0

Boolean

A truth value with two possible values.

Samples:
true
false

Float

A floating point number with the range 1.40129846432481707e-45 to 3.40282346638528860e+38 (positive and negative).

Samples:
4.6
-2.0

An **Integer** is automatically converted into a Float, if it is necessary.

String

A string of characters. Constant strings can either be written in quotation marks or in tick marks. When using quotation marks, tick marks can be written inside the string and vice versa without escaping.
See **Escape characters** to get a list of supported escape characters.

Samples:
'word'
"long text"
"A string with\n\t'tick marks' and\n\t\"quotation marks\""

Character

A single character.
See **Escape characters** to get a list of supported escape characters.

Samples:
'A'
'\n'

Escape characters:
Within constant **String**s or **Character**s, the following expressions are valid escape characters:

\n	Line break
\t	Tab
\\	Backslash
\"	Quotation marks
\'	Tick mark

4.2 Graph types

Graph

A graph including vertices, edges, configuration, vertex- and edge data, name and algorithm cache.

Sample: getCurrentGraph()

Vertex

A single vertex of a graph. A constant vertex can be written with `$[vertex identifier]`. This will deliver the vertex with the given identifier of the respective graph. If no such vertex exists, an error will occur.

Samples:
$v2
getVertexByIdent('v2')

Edge

A single edge of a graph. An edge can be identified by the two incident vertices (in directed graphs by their ordering, too).
An edge can for example be delivered using *vertex1* `--` *vertex2* in undirected and *vertex1* `->` *vertex2* in directed graphs.

Samples:
getEdge($v0,$v2)
$v0 -- $v2

4.3 Complex types

Vector

A vector holds multiple ordered values and can be of any size. An entry of the vector can be of any individual type (also Vector again). The particular entries can be accessed with *vector*[*index*] where *index* is an **Integer** beginning at `0`.
Vectors can be used in foreach-loops (see 'graphana_manual.pdf').

Samples:
(1,2,3,4)
()
((2.4,4.2,6.4),(-5.6,7,10.2),(3,2.1,0))

See **Complex type operations** for a list of operations on vectors.

Set

A set holds multiple unordered values and can be of any size. An element of the set can be of any individual type (also Set again).
Sets can be used in foreach-loops (see 'graphana_manual.pdf').

Samples:
{1,2,3,4}
{}
{"A string",$aVertex,{2.3,8.5,-6}}

See **Complex type operations** for a list of operations on sets.

Iterable

An Iterable cannot be created directly. An argument of an operation call is casted into an Iterable if it is a **Vector** or a **Set**.

4.4 File types

File

Files are given as **String**s containing the relative or absolute filename. The file does not have to exist.

Samples:
”C:/absolute/path/file.ext”
”relative/path/file.ext”

ExistingFile

Nearly similar to **File** but the file must exist.

Sample: ”path/to/an/existing/file.ext”

4.5 Miscellaneous types

Histogram

A histogram contains a mapping from integral numbers to float numbers. Every entry can be accessed in particular.

Samples:
createHistogram(30)
degreeDistribution()

See **Histogram creation** for a list of operations on histograms.

Interval

An interval has a minimum and a maximum value (”bounds”). The bounds can be extracted using **getLowerBound** and **getUpperBound**

Samples:
newInterval(-3,7)
cvdBounds()

ParseTree

A **String** can be parsed and converted into a parse tree. This tree can then be executed any number of times and does not need to be parsed again, which improves performance. Note that the execution of a tree may return different results depending on global variables for example. So if a script shall be executed very often, it makes sense to convert it into a parse tree once using **parse**, assign it to a variable and execute it repeatedly (for example within a loop) using **executeTree**.

Sample: parse(”2 + x*(3+y)”)